

Random graphs, social networks and the internet

Lecture 1

Mariana Olvera-Cravioto

UNC Chapel Hill

`molvera@unc.edu`

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Social networks and graphs

- ▶ The internet, the web, Facebook, Twitter, LinkedIn, Instagram, WhatsApp, WeChat, Snapchat, Pinterest, Reddit, etc. are all examples of **networks**.
- ▶ In **social networks**, connections occur among people.
- ▶ A connection between two people can mean many different things depending on the network, e.g., friendship, hyperlinks, follower-followed relations, etc.
- ▶ There are also many networks that do not involve people at all, e.g., the internet, neural connections in the brain, interactions between proteins in biology, articles in a citation network, etc.
- ▶ When analyzing networks, it is often convenient to think of them as **graphs**.

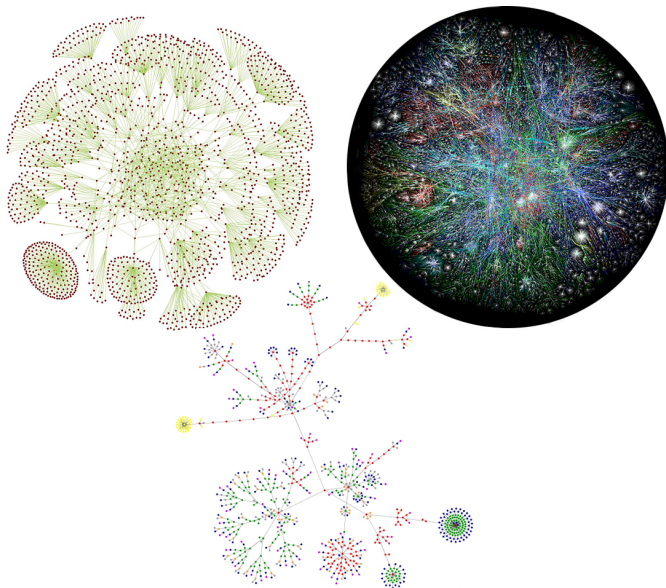
Graphs

- ▶ A graph consists on a set of **vertices**, V , and a set of **edges** E .
- ▶ Graphs can be **undirected** or **directed**.
- ▶ In an undirected graph, the relation between the vertices is symmetric, while in a directed graph it is not.
- ▶ We will call the vertices $V = \{1, 2, \dots, n\}$, and write $i \rightarrow j$ to mean there is an edge (perhaps undirected) from vertex i to vertex j .
- ▶ In an undirected graph, the **degree** of a vertex is the number of edges incident to it.
- ▶ In a directed graph, the **in-degree** is the number of inbound edges and the **out-degree** is the number of outbound edges.

Examples of graphs

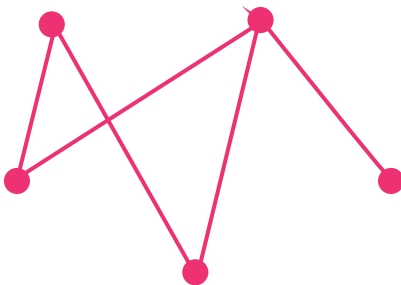
- ▶ **Internet:** vertices are servers/computers and edges are connections between them.
- ▶ **Web:** vertices are webpages and edges are directed links from one page to another.
- ▶ **Facebook:** vertices are people and edges are friendship relations.
- ▶ **Twitter:** vertices are people and edges go from followers to followed.
- ▶ **Citation network:** vertices are research articles and a directed edge represents a citation.

Different types of graphs



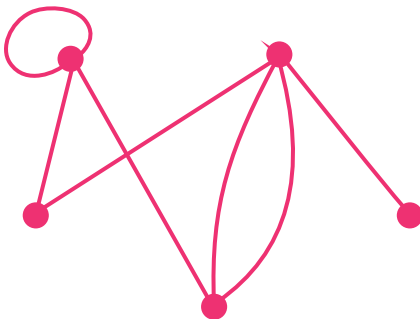
Types of graphs

- **Simple graphs:** a graph that has no self-loops nor multiple edges between any two vertices.



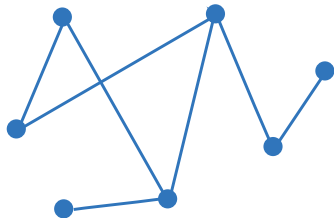
Types of graphs

- **Multigraphs:** a graph that may have self-loops or multiple edges between two vertices.

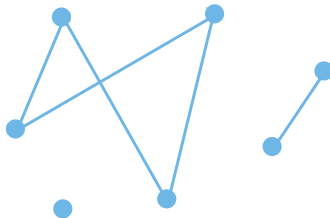


Types of graphs

- **Connected graphs:** graphs where every pair of vertices is connected through a path.



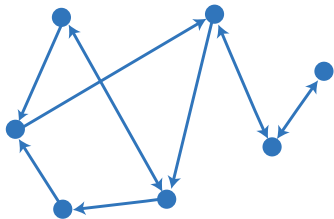
Connected



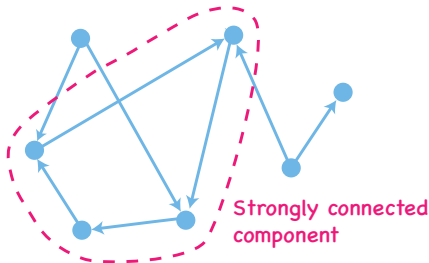
Disconnected

Types of graphs

- **Strongly connected graphs:** for directed graphs and any pair of vertices i and j , there exists a directed path from i to j and one from j to i .



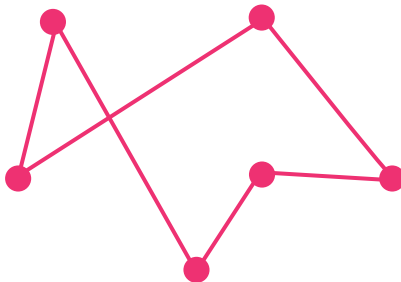
Strongly connected



Weakly connected

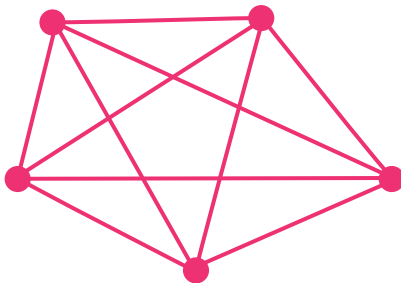
Types of graphs

- **Regular graphs:** all the vertices in the graph have the same degree.



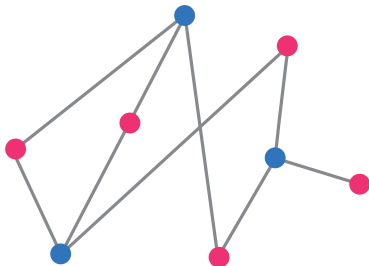
Types of graphs

- **Complete graphs:** there is an edge between every pair of vertices in the graph.



Types of graphs

- **Bipartite graphs:** there are two classes of vertices, say V_1 and V_2 , and edges occur only between a vertex in V_1 and one in V_2 .



Structures and properties

- ▶ Some structures that can be of interest when studying graphs are:
 - ▶ **Cycles:** paths that start and end with the same vertex without repeated vertices.
 - ▶ **Cliques:** complete subgraphs.
 - ▶ **Distance between two vertices:** length of the minimum path connecting two vertices; in directed graphs the path must be directed.
 - ▶ **Component of a vertex:** the set of vertices that can be reached through (directed) paths from a given vertex.
- ▶ Some properties of interest:
 - ▶ **Diameter:** the maximum distance between two points in the graph.
 - ▶ **Components:** sizes of the largest, second largest, etc.
 - ▶ **Cycle lengths:** the typical length of cycles in the graph.
 - ▶ **Clustering:** the proportion of triangles (3-cliques) vs. open wedges.
 - ▶ **Communities:** subsets of vertices that have more edges among their vertices than with vertices outside the set.

Some questions of interest

- ▶ Is the graph (strongly) connected?
 - ▶ If not, does there exist a giant (strongly) connected component?
 - ▶ What is the size of the smaller components?
- ▶ What is the diameter of the graph?
- ▶ What is the typical distance between vertices in the graph?
- ▶ What is the degree distribution in the graph?
- ▶ Does the graph have clusters/communities?
- ▶ Are there vertices that are more “influential” or “central” to the network?

The small world phenomenon

- ▶ In the late 60's, a social psychologist named Stanley Milgram conducted a set of experiments to try to determine the typical length of paths connecting two individuals in the United States.
- ▶ A letter addressed to somebody in Boston would be given to a set of randomly chosen people in different states in the Midwest, strangers to the recipient, with the instruction to help it reach its destination by sending it to an acquaintance.
- ▶ **Result:** it took an average of 6 people to connect the first sender and the final recipient, something that became known as the
small world or six degrees of separation
phenomenon.
- ▶ Interestingly, the small world property is very common in large real-world networks.

Scale-free networks

- ▶ Recall that the degree of a vertex $i \in V = \{1, 2, \dots, n\}$ in an undirected graph, denoted D_i , is the number of edges incident to it.
- ▶ The proportion of vertices having degree $k = 0, 1, 2, \dots$, is given by

$$p(k) = \frac{1}{n} \sum_{i=1}^n 1(D_i = k)$$

- ▶ We call $\{p(k) : k \geq 0\}$ the **degree distribution**.
- ▶ If the degree distribution of a graph satisfies

$$p(k) \propto k^{-\gamma}$$

for some $\gamma > 0$ (usually $\gamma \in (2, 3)$), we say that the graph is **scale-free**.

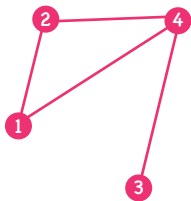
- ▶ In a scale-free graph there are vertices that have really large degrees, even if the average degree is small.

The adjacency matrix

- ▶ A convenient way to represent a graph is through its adjacency matrix.
- ▶ For a graph $G(V, E)$ having vertices $V = \{1, 2, \dots, n\}$, its adjacency matrix A is the $n \times n$ matrix which has:

$$a_{i,j} = \text{number of edges from vertex } i \text{ to vertex } j$$

- ▶ In a **simple** graph we have $a_{i,j} \in \{0, 1\}$ for all (i, j) .
- ▶ In a **undirected** graph the matrix A is symmetric.
- ▶ Example:



$$A = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

Counting paths

- ▶ Suppose A is the adjacency matrix of a graph $G(V, E)$.
- ▶ Then, for any $k \geq 1$ and any $i, j \in V$,

$$(A^k)_{i,j}$$

is the **number of paths of length k** from vertex i to vertex j , where A^k is the matrix A raised to the k th power.

Random graph models

- ▶ Some real networks are too big to be analyzed exactly.
- ▶ Some may even be constantly changing.
- ▶ **Idea:** we can think of our specific real-world graph as just one “typical” element of a larger class.
- ▶ If we can show that a property holds for a large class of graphs, it is likely it will hold for our specific graph.
- ▶ **Random graphs** are mathematical models that can help us understand large real-world graphs.
- ▶ No random graph model can mimic all the properties of a specific real-world graph, so we focus on choosing models that share certain properties that are important to the problem we want to analyze.

Large graph limit

- ▶ Random graph models consist of a vertex set $V_n = \{1, 2, \dots, n\}$ and a set of rules for determining whether a given edge is present or not based on some random events.
- ▶ Their mathematical analysis is usually done under the **large graph limit** $n \rightarrow \infty$ on a sequence of graphs $\{G(V_n, E_n) : n \geq 1\}$.
- ▶ Taking the limit $n \rightarrow \infty$ simplifies computations in order for us to identify general properties.
- ▶ In practice, establishing results in the large graph limit means that our findings are likely to be true for sufficiently large graphs.

Static vs. evolving models

- ▶ Random graph models can be broadly classified into two categories: **static models** and **evolving or growing models**.
- ▶ Static models are meant to represent a “snapshot” of a large network.
- ▶ In static models $G(V_n, E_n)$ and $G(V_{n+1}, E_{n+1})$ can be totally different.
- ▶ Evolving models are meant to describe the growth of a graph as vertices get added to the graph (usually one at a time), so $G(V_n, E_n)$ and $G(V_{n+1}, E_{n+1})$ share most edges.
- ▶ In many evolving models edges and vertices never disappear, so $G(V_n, E_n)$ is a subgraph of $G(V_{n+1}, E_{n+1})$.

The Erdős-Rényi random graph

- ▶ The simplest model for a random graph is the **Erdős-Rényi model**.
- ▶ Consider a graph with vertex set $V_n = \{1, 2, \dots, n\}$.
- ▶ There are a total of $\binom{n}{2}$ possible edges in the graph, and each of them will be chosen to be present or not with a coin flip.
- ▶ Suppose you have a coin that lands heads with probability $p \in (0, 1)$.
- ▶ For each pair of vertices i and j , toss the coin; if it lands heads, draw an edge between i and j , otherwise do nothing.
- ▶ Equivalently, if A denotes the adjacency matrix of the graph, let

$$a_{i,j} = a_{j,i} = 1(\text{coin-flip is a head}), \quad i \neq j,$$

and set $a_{i,i} = 0$.

Properties of the Erdős-Rényi model

- ▶ This is the most studied random graph model there is.
- ▶ Some of its connectivity properties are:
 - ▶ If $np < 1$ the graph will consist of only small components of size $O(\log n)$.
 - ▶ If $np \rightarrow c > 1$ the graph will contain a unique *giant* connected component, with all other components of size $O(\log n)$.
 - ▶ If $np = 1$ the largest component will have size $O(n^{2/3})$.
 - ▶ If $p < (1 - \epsilon)n^{-1} \log n$ the graph will most likely be **disconnected**.
 - ▶ If $p > (1 + \epsilon)n^{-1} \log n$ the graph will most likely be **connected**.
- ▶ When the graph is connected, it exhibits the **small-world** property, with typical distance of order $O(\log n)$.

Degree distribution

- ▶ To compute the degree distribution we can use binomial probabilities.
- ▶ Fix a vertex $i \in V_n$, then its degree is given by

$$D_i = \sum_{j=1}^n \chi_{i,j}, \quad \chi_{i,j} = 1((i,j) \in E_n)$$

- ▶ Note that the $\chi_{i,j}$ are independent Bernoulli r.v.s with parameter p .
- ▶ Therefore, since all vertices have the same distribution, for all $i \in V_n$,

$$P(D_i = k) = P(D_1 = k) = P(\text{Binomial}(n, p) = k) = \binom{n}{k} p^k (1-p)^{n-k}$$

- ▶ Moreover, if $np \rightarrow c$ as $n \rightarrow \infty$, we have that

$$\lim_{n \rightarrow \infty} P(D_1 = k) = \frac{e^{-c} c^k}{k!}, \quad k \geq 0,$$

known as a Poisson distribution with mean c not **scale-free**.

Thank you for your attention.